

Hardware & Software Verification

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Lecture 6

Last lecture

- We need to be able to **reason about** the programs we write, not merely **test** them. There is a large and growing need for this.
- Dafny is a **verification-oriented** programming language. Its compiler will refuse to produce executable code until it has proven the code to be **correct**.

But what does
correct mean?

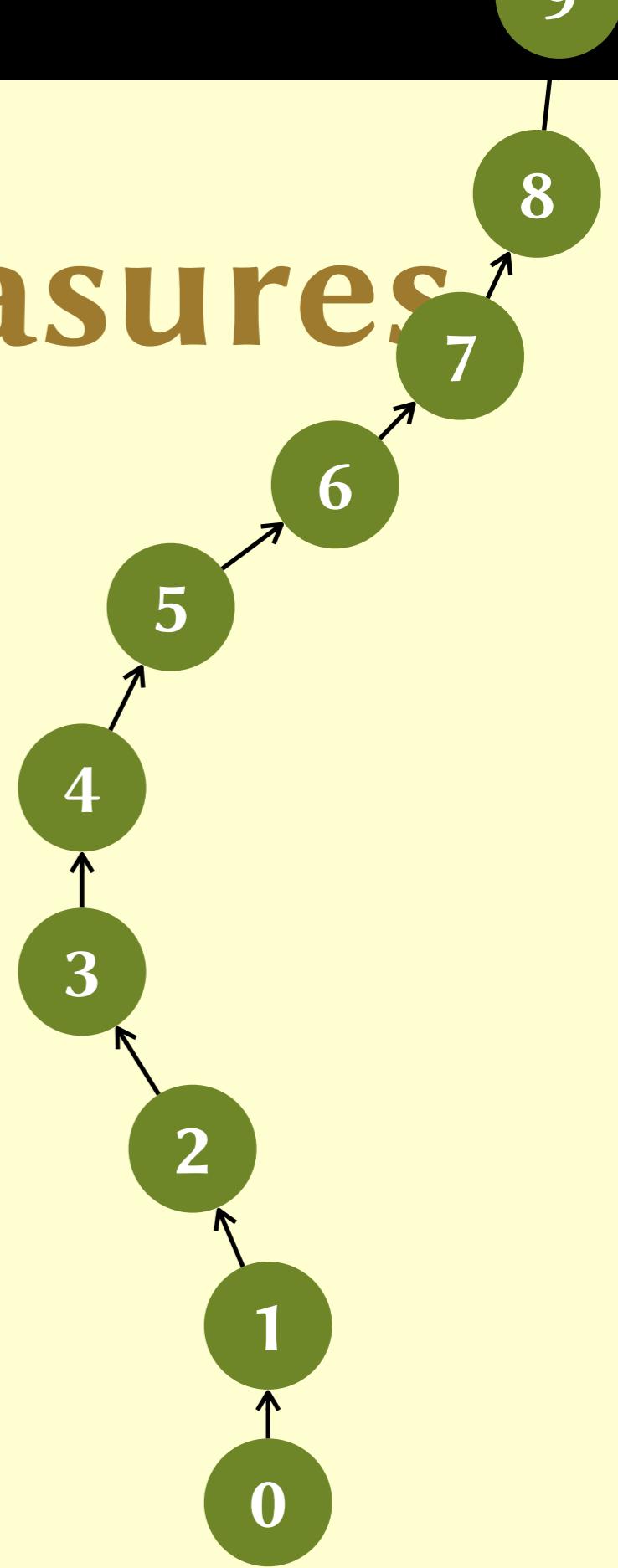
Demo: max of a pair

- named output parameters
- postconditions
- overly weak/strong specifications

Demo: max of an array

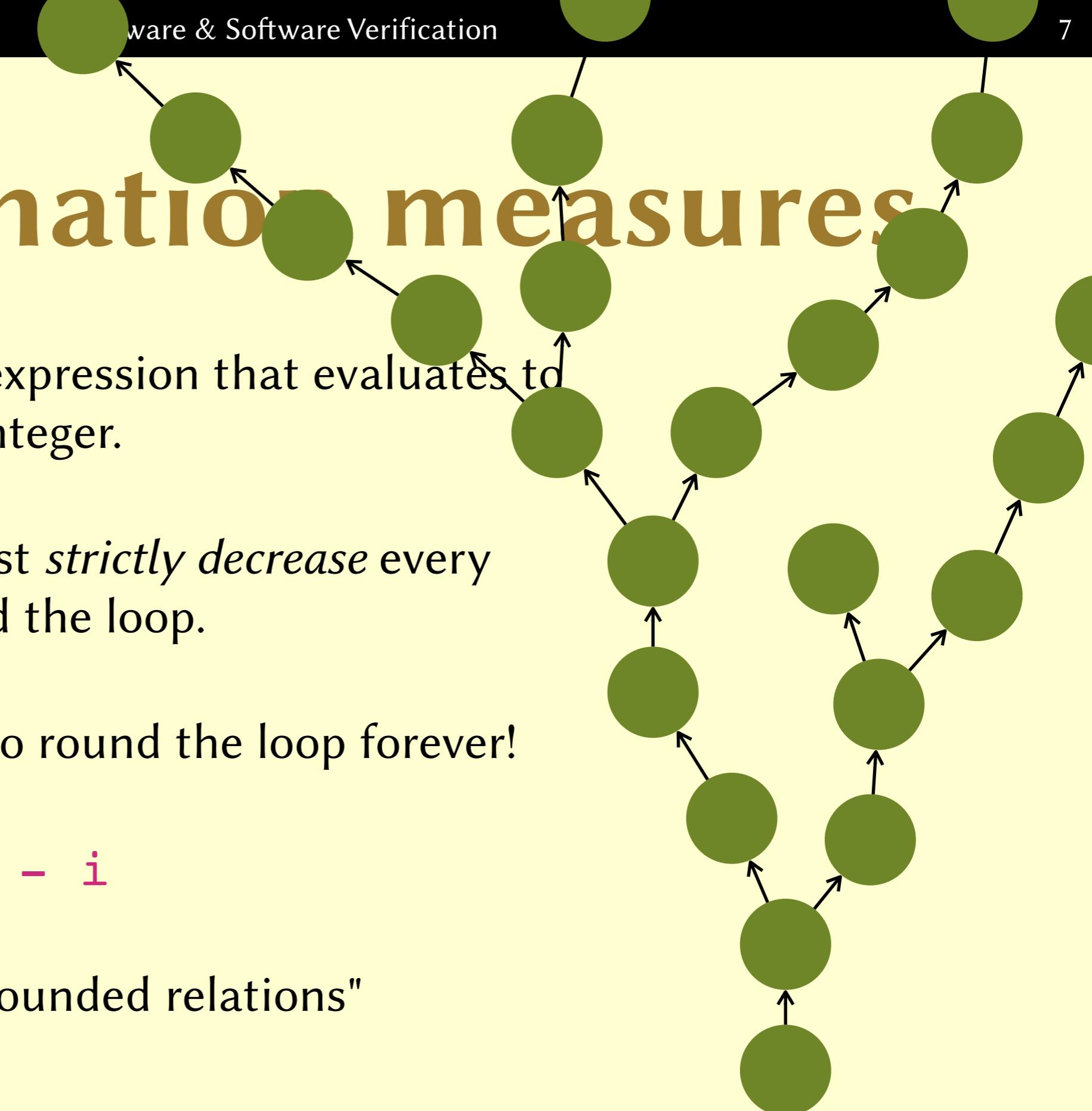
Termination measures

- A *measure* is an expression that evaluates to a non-negative integer.
- The measure must *strictly decrease* every time we go round the loop.
- Hence we can't go round the loop forever!
- E.g.: `A.Length - i`
- "Theory of well-founded relations"



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- Hence we can't go round the loop forever!
- E.g.: **A.Length - i**
- "Theory of well-founded relations"



Demo: max of an array

The problem with loops

code
before
loop

invariant

postcondition?

code
before
loop

invariant
body
invariant
postcondition?

code
before
loop

invariant

body

invariant

body

invariant

postcondition?

code
before
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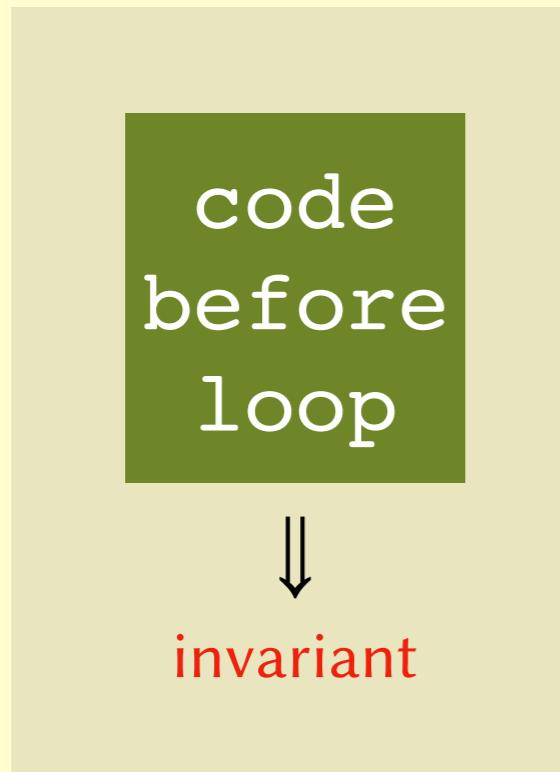
invariant
body
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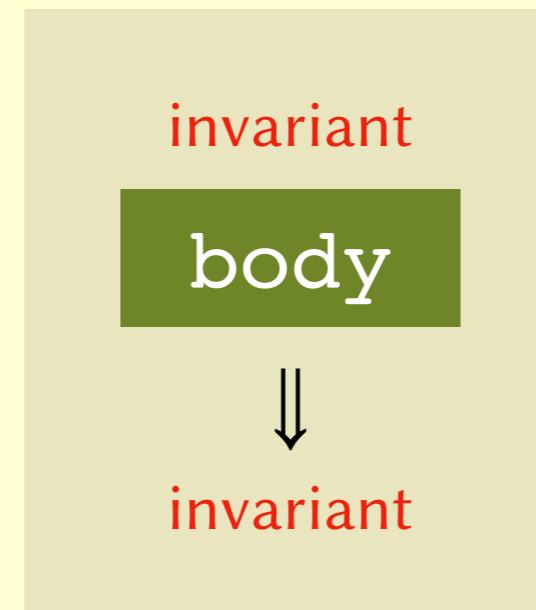
invariant
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Loop invariants

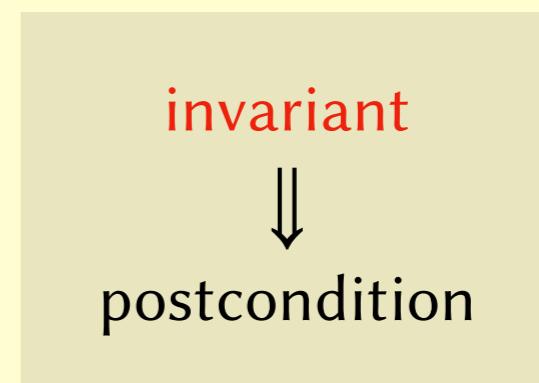
1.



2.



3.

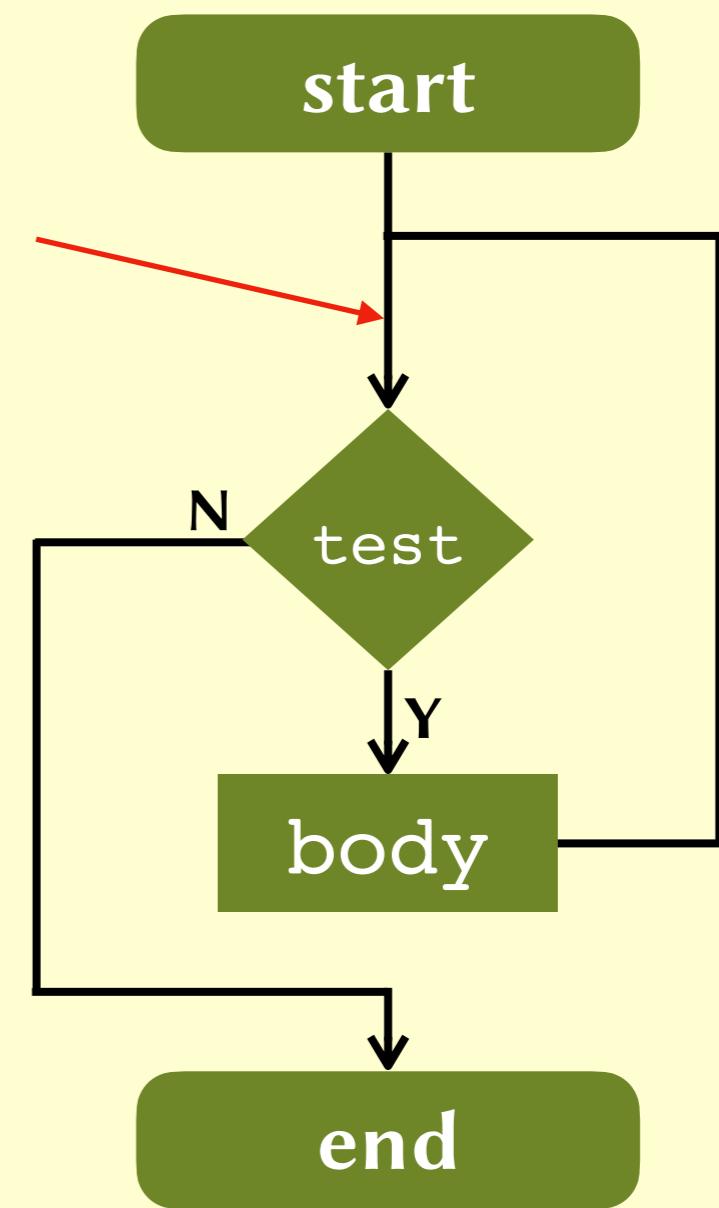


Demo: max of an array

Loop invariants

```
while test
  invariant foo
{
  body
}
```

foo must
hold here!



Finding invariants

A[0]	A[1]	A[2]	A[3]	A[4]	A[5]	A[6]
4	0	1	9	7	1	2

```
r := A[0];
var i := 1;
while i < A.Length {
  if r < A[i] {
    r := A[i];
  }
  i := i+1;
}
```

i	r
1	4
2	4
3	4
4	9
5	9
6	9
7	9

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i	r	$\exists j. 0 \leq j < i \wedge r = A[j]$
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Demo: max of an array

- syntax for variables (**var**) and arrays (**array<...>**)
- preconditions (**requires**)
- termination measures (**decreases**)
- universal (**forall**) and existential (**exists**) quantification
- loop invariants (**invariant**)
- predicates (**predicate**)

Coursework 1

- Worksheet is now on Github!
- All coursework is due **Friday 16th December at 23:59.**
- Please submit a single Dafny source file via Teams.
- Please include lots of */*comments*/* in your source file to explain your thinking.
- Please work in pairs, and submit one file per pair.
- Please do not share your answers with other pairs.
- If you have questions, please come to the Monday labs or raise an issue on Github.